

GLAST Large Area Telescope

Instrument Flight Software

**Inter-Task Communication
Functional Demonstration
April 15, 2005**

Stanford Linear Accelerator Center



Structure of the Inter-task Communications System

- The integrated inter-task communications system is built from three packages:
 - The ITC package defines the Inter-Task Communications standard:
 - ITC defines a common communications standard to unify communications between tasks on either the same or different CPUs.
 - It also provides the services to build up tasks that are capable of communicating according to the standard.
 - ITC is not itself a task, it's a toolbox.
 - The CTS package implements the ITC communications standard on 1553 transport hardware
 - It provides the service layer between applications (high level tasks) and the Command and Telemetry Database (CTDB) hardware driver.
 - It is constructed using the ITC toolbox, allowing CTS to present an ITC compliant (uniform) interface to the application layer tasks.
 - It implements tasks to communicate:
 - Spacecraft to CPU
 - CPU to spacecraft



Structure of the Inter-task Communications System (2)

- The LCS package implements the ITC communications standard on LCB/LATp transport hardware:
 - It provides the service layer between applications (high level tasks) and the LCB hardware driver.
 - It is constructed using the ITC toolbox, allowing LCS to present an ITC compliant (uniform) interface to the application layer tasks.
 - It implements tasks to communicate:
 - CPU to CPU
 - CPU to SSR

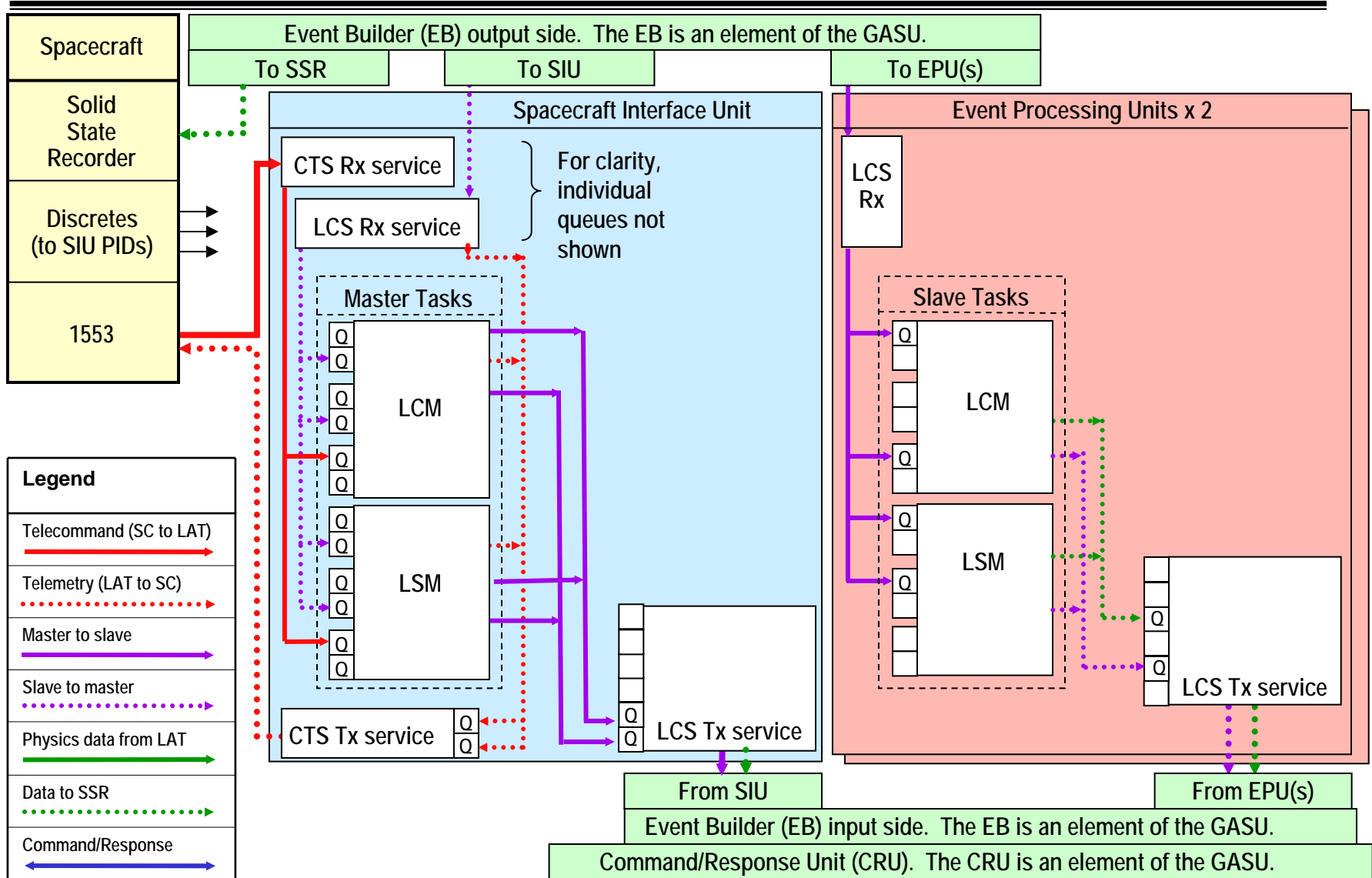


Dimensions Exercised

- The ITC framework defines the following dimensions of inter-task communications:
 - Nodes: up to 8 supported -- includes CPUs, Spacecraft, SSR, broadcast
 - Tasks: 32 tasks maximum -- applications, services, and “lightweight tasks”
 - Queues:
 - Up to 256 queues per CPU
 - Maximum of 8 queues per task
 - Highest (META_HI) and lowest (META_LO) queues reserved to ITC. Remaining six queues available to application designer.
 - Communications protocols: dispatch and hardware
 - Dispatch protocols: RAW, CMD, APID
 - Hardware protocols for LATp fabric: NOACK, ACK
- For the demonstration:
 - Three nodes: SIU (node 0), EPU0 (node 1), EPU1 (node 2)
 - Two tasks: LCM (task 1), LSM (task 8)
 - Two queues: CTL_X (queue 4), CTL (queue 5)
 - Two hardware protocols: NOACK (protocol 2), ACK (protocol 3)
 - There will also be a SC (node 5), but the demonstration is focused on LAT internal communications (package LCS)



System Context for the Demo



15 April 2005

FSW Functional Demonstration

5



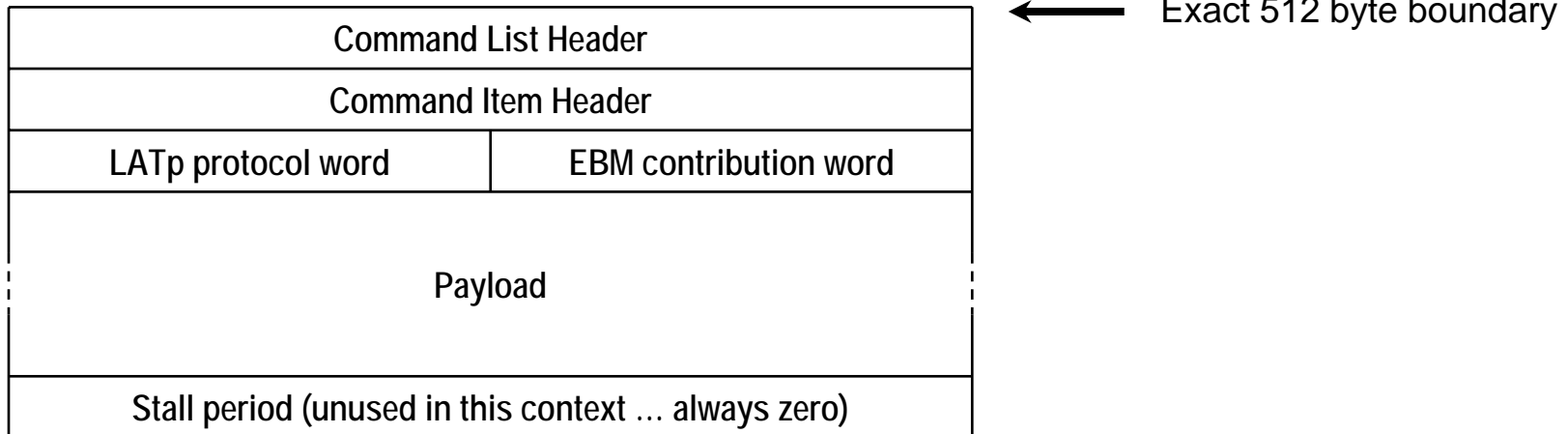
Communications Protocols Used by LCS

- LCS supports 2 communications protocols
 - Protocol 2 (NOACK): Unacknowledged (“fire and forget”)
 - Protocol 3 (ACK): Acknowledged
- Protocol 2
 - Uses
 - Talk to EPU in primary boot (primary boot code cannot acknowledge)
 - Talk to SSR (SSR cannot acknowledge)
 - CPU broadcasts
 - Characteristics
 - Messages limited to a single packet (2 kByte limit)
 - Except for SSR (special case)
 - Completion notification comes when message has “left the building”
 - Hence, no timeout behavior
- Protocol 3
 - Uses
 - General CPU to CPU communications
 - Characteristics
 - Messages limited to 128 MByte (i.e. no effective limit)
 - Completion notification comes from delivery of message to receiving task/queue on *destination* CPU
 - Full timeout behavior

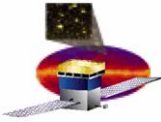


Notes on LCB Hardware Constraints

- LCB hardware places quite onerous constraints on requests it will service
 - The memory layout must be very precise



- LCB hardware limits a single request to (4096 – 128) bytes
- Among other services, LCS makes these constraints invisible to the applications writer
 - Messages are realigned correctly before transmission
 - Long messages are broken down into a series of packets on the sender and reassembled into the original message on the receiver



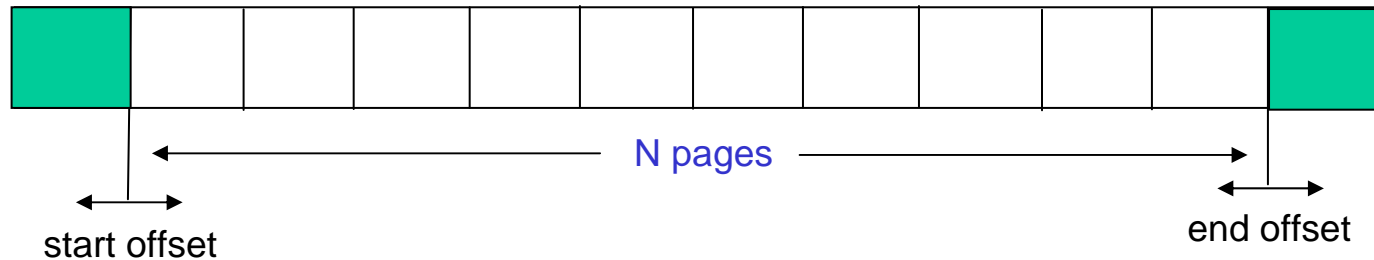
Sending A Test Message (1)

- The test suite must synthesize messages covering a wide parameter space
 - Both hardware protocols
 - A range of lengths
 - A range of alignments
 - Verification of packet sent against packet received
 - Verification of NON-overwrite of “un-owned” memory
- The messages can then be sent from any source to any destination
- All the tests demonstrated today use the same message description
 - A base of N pages (a page is 512 bytes)
 - A byte offset (positive or negative) on the start of message
 - A byte offset (positive or negative) on the end of message



Sending A Test Message (2)

- What happens
 - (N+2) pages of memory are allocated
 - The buffer is painted with predictable pattern



- What gets sent
 - The section of the buffer between the requested offsets

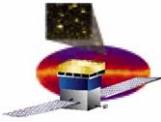


- When message transmission completes
 - Check that buffer still contains the predicted pattern
 - Free the buffer



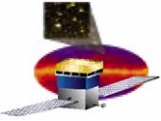
Receiving A Test Message

- When message first received
 - Allocate a receive buffer (in fact it's slightly over-allocated)
 - Prepaint buffer with fixed pattern (`0xdeafbede`)
 - Overwrite buffer with incoming message
- When message completes (may be many packets)
 - Forward to destination task/queue
- At receiving task/queue
 - Check that buffer contains predicted pattern (NOT `0xdeafbede`)
 - Check that area *past* the message *does* contain `0xdeafbede`
 - Free the buffer
- Every message sent today will be subjected to these verifications
 - Any flaw will result in a printed error message
 - Exceptions for raw performance tests and deliberately broken messages

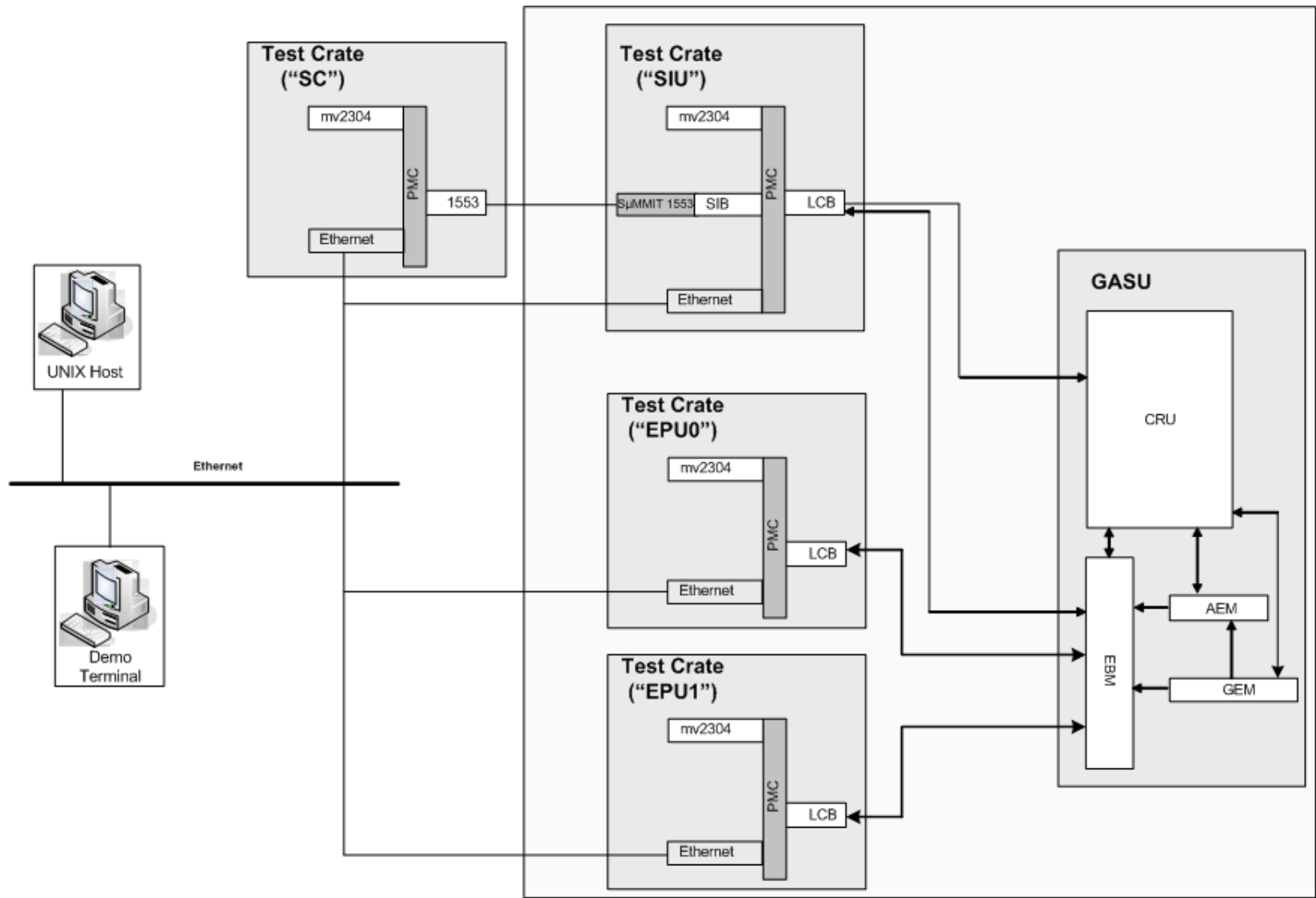


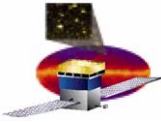
Demonstration Technique

- Terminal windows will be opened to monitor the Spacecraft crate and CPU crates used in the demo
 - Routines from the LCS test suite will be invoked from the command line to exercise various pieces of the inter-task communications system
 - Results will be displayed at the terminal windows



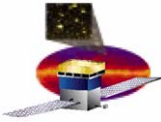
Hardware Context





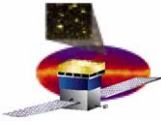
On to the Demonstrations

- Send single message from arbitrary source node/task to arbitrary destination node/task/queue with arbitrary protocol
 - Audience participation invited
- Pratfall tests
 - Bad destination node
 - Bad destination task
 - Bad destination queue
- Speed test
 - Send 4 large messages back-to-back and time them
- Priority test
 - Send 1 short, high priority message through long, multi-part low priority message
- Send a long series of messages using protocol 2 (NOACK)
 - Takes about 1 minute to complete



And More Demonstrations

- Corruption tests
 - Use low-level LCB code to write errors into the outbound data, mimicking a hardware corruption:
 - Broken sequence number
 - Broken transaction ID
 - Broken length
 - Stalled transmission
- Torture test (the hardware, not me)
 - Audience participation invited to select:
 - 4 sources on SIU to 4 destinations on EPU0 or EPU1
 - 4 sources on EPU0 to 4 destinations on SIU or EPU2
 - 4 sources on EPU1 to 4 destinations on SIU or EPU1
 - Transmit a long series of messages for each combination of source and destination (about 15 minutes)
 - Review statistics for this large block of exchanged messages
- Maybe a surprise demo thrown in!



Forward Work

- Complete the test suite
 - Ensure that reconstruction of *hardware* fragmentation works
 - Ensure all timeout race conditions are accounted for
- Implement the CPU -> SSR service
 - Easier
 - SSR can't talk back! (no protocol 3)
 - Only one queue (no point and *dangerous* to interleave messages)
 - Harder
 - Packetization must follow CCSDS rules (14 byte headers ... aarrghh)
 - Left to last because
 - No immediate customers
 - Only clumsy test methods available (bring on the VSC)
- Implement broadcast
 - May be omitted (no identified customers)
- Write the manual